Two bits were received from each mobile station, which must be decoded separately.

$$\begin{aligned} R_{S1}^{1} &= K_{S1} \cdot R_{BS}^{1} = (+1, +1, -1, +1) \cdot (0, 0, +3, -2) = -5 \rightarrow "0" \\ R_{S1}^{2} &= K_{S1} \cdot R_{BS}^{2} (+1, +1, -1, +1) \cdot (+2, 0, -2, +2) = +6 \rightarrow "1" \\ R_{S2}^{1} &= K_{S2} \cdot R_{BS}^{1} = (+1, +1, +1, -1) \cdot (0, 0, +3, -2) = +5 \rightarrow "1" \\ R_{S2}^{2} &= K_{S2} \cdot R_{BS}^{2} (+1, +1, +1, -1) \cdot (+2, 0, -2, +2) = -2 \rightarrow "0" \end{aligned}$$

b)

Without noise:

$$\begin{split} R_{BS}^{1} &= D_{S1}^{1} \times K_{S1} + D_{S2}^{1} \times K_{S2} \\ &= (-1) \times (+1, +1, -1, +1) + (+1) \times (+1, +1, +1, -1) \\ &= (0, 0, +2, -2) \\ R_{BS}^{1} &= D_{S1}^{2} \times K_{S1} + D_{S2}^{2} \times K_{S2} \\ &= (+1) \times (+1, +1, -1, +1) + (-1) \times (+1, +1, +1, -1) \\ &= (0, 0, -2, +2) \\ R_{BS} &= R_{BS}^{1} |R_{BS}^{2} = (0, 0, +2, -2, 0, 0, -2, +2) \end{split}$$

Now we must only subtract from the actually received sequence in order to get the noise:

$$N = (0,0,+3,-2,+2,0,-2,+2) - (0,0,+2,-2,0,0,-2,+2)$$

= (0,0,+1,0,+2,0,0,0)

c)

In order to present good characteristics against multipath fading, a key must be significantly orthogonal to rotated versions of itelf. Regarding K_{S1} :

$$K_{S1} \cdot (K_{S1} \gg 1) = (+1, +1, -1, +1) \cdot (+1, +1, +1, -1) = 0$$

$$K_{S1} \cdot (K_{S1} \gg 2) = (+1, +1, -1, +1) \cdot (-1, +1, +1, +1) = 0$$

$$K_{S1} \cdot (K_{S1} \gg 3) = (+1, +1, -1, +1) \cdot (+1, -1, +1, +1) = 0$$

As such, signals from S1 can endure a delay spread up to 3 chip times. Similar results can be obtained for K_{S2} .

1.

a)

Single hop:

$$P_t^{1-hop}[dBm] = P_r[dBm] + PL_0 - G_t[dBi] - G_r[dBi] + 10 \cdot \alpha \cdot \log_{10}\left(\frac{2d}{d_0}\right)$$

$$P_t^{1-hop} = \frac{P_r \cdot 10^{\frac{PL_0}{10}} \cdot \left(\frac{2d}{d_0}\right)^{\alpha}}{G_t \cdot G_r} = \frac{P_r \cdot 10^{\frac{PL_0}{10}} \cdot (2d)^3}{G_t \cdot G_r}$$

With two hops we have, for each hop:

$$P_t^{2-hop} = \frac{P_r \cdot 10^{\frac{P_{L_0}}{10}} \cdot d^3}{G_t \cdot G_r}$$

We can now divide the power used for one hop by the power used for two hops:

$$\frac{P_t^{1-hop}}{2 \cdot P_t^{2-hop}} = \frac{\frac{P_r \cdot 10^{\frac{PL_0}{10}} \cdot 2^3 \cdot d^3}{G_t \cdot G_r}}{2\frac{P_r \cdot 10^{\frac{PL_0}{10}} \cdot d^3}{G_t \cdot G_r}} = \frac{2^3 \cdot d^3}{2 \cdot d^3} = 4 > 1$$

b)

There will be two concurrent signals at C, plus termal noise. The signal from B constitutes the useful signal, while the signal from A will constitute interference. The first step is to calculate the power received for each of these signals. Since the transmit power corresponds to the minimum, the received power from B is equal to the receiver sensitivity: -80 dBm. We only have to calculate the received power from A. We know that at 100 m from A, the power is -80 dBm. We can use this reference distance and power and proceed to calculate the received power at C, which is located at 200 m from A:

$$P_r^A[dBm] = P_t [dBm] - PL_0 - 10 \cdot \alpha \cdot \log_{10} \left(\frac{d}{d_0}\right) = -80 - 10 \cdot 3 \cdot \log_{10} \left(\frac{200}{100}\right)$$

\$\approx -89.0 dBm\$

The interference plus noise power is thus:

$$I + N = 10^{\frac{-89}{10}} + 2000000 \times 10^{\frac{-20}{10}} \, mW \approx -89.0 \, dBm$$

The SINR is then:

$$P_r^B[dBm] - (I+N)[dBm] \approx 9.0 \ dB = 8.0$$

c)

We must first calculate the BER, which requires the *Q* function. Knowing that R = B (since r = 0):

$$BER_{BPSK} = Q\left(\sqrt{\frac{2 \cdot E_b}{N_0}}\right) = Q\left(\sqrt{\frac{2 \cdot SINR \cdot B}{R}}\right) = Q\left(\sqrt{\frac{2 \cdot 8.0 \cdot 2000000}{2000000}}\right)$$
$$\approx 3.17 \times 10^{-5}$$

Now, we can calculate the FER:

$$FER = 1 - (1 - BER_{BPSK})^{20 \times 8} \approx 4.94 \times 10^{-2}$$

Packet size
$$(l) = 64 \text{ kbps} \times 40 \text{ ms} = 2560 \text{ bits} = 320 \text{ bytes}$$

$$T_{packet} = DIFS + Backoff + PHo + \frac{(MACh + RTP/UDP/IP + data)}{R} + SIFS + PHo + \frac{ACK}{R} =$$

= 0.034 + 0.067 + (0.096 + $\frac{34 \times 8 + 40 \times 8 + 320 \times 8}{2000}$) + 0.016 + 0.096 + $\frac{14 \times 8}{2000} \approx 1.94 \text{ ms}$
 320×8

Throughput =
$$\frac{320 \times 8}{1.94}$$
 kbps ≈ 1320 kbit/s

b)

 $N = [1320/64] = 21 \text{ streams} \rightarrow \text{Since each call is bi-directional} \rightarrow 10$ telephones

c)

$$FER_{MAC} = (FER_{PHY})^{1+4} = 0.03^5 = 2.43 \times 10^{-8}$$

4.

a)

$$\frac{R}{W} = 4 \Leftrightarrow W = \frac{R}{4} \Leftrightarrow W = \frac{12000000}{4} = 3 MHz$$

b)

The provided bitrate value is a diversion, since we can always trade-off bandwidth for capacity. We just have to apply the Shannon-Heartley theorem, taking into account that $\frac{c}{W} \approx 10$ for $\frac{E_b}{N_0} = 20$:

$$C = W \cdot (1 + SNR) \Leftrightarrow \frac{C}{W} = (1 + SNR) \Leftrightarrow SNR = 2^{10} - 1 = 1023$$

c)

1 symbol = 2 bits

$$1 - P_M = (1 - BER)^2 \Leftrightarrow BER \approx 5 \times 10^{-6}$$

3. a)