

Multitasking Operating Systems

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2015-10-12

- Motivation for processes
- The process abstraction
- Context switching
- Multitasking

- Respond to external events
 - Engine controller
 - Seat belt monitor

- Real-time response has implications
 - System architecture
 - Program implementation

- May require a chain reaction among multiple processors

Processes help us manage timing complexities

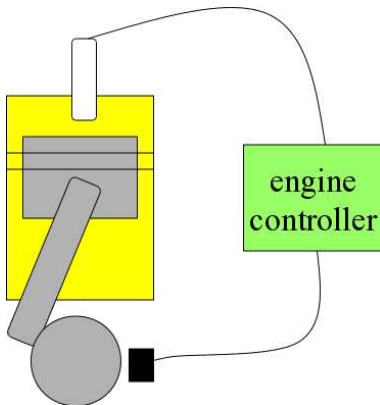
- Multiple rates
 - multimedia
 - automotive

- Asynchronous input
 - user interfaces
 - communication systems

- Tasks may be synchronous or asynchronous
- Synchronous tasks may recur at different rates
- Processes run at different rates based on computational needs of the tasks

Period of a process

- Is the time between successive executions
- Is the inverse of the **rate**



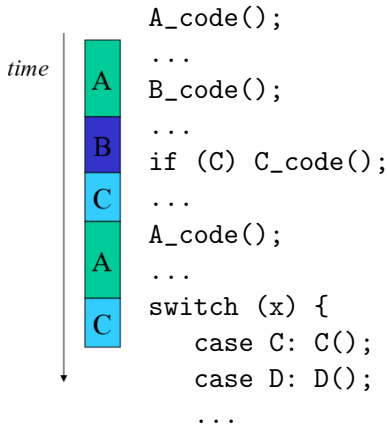
Tasks

- spark control
- crankshaft sensing
- fuel/air mixture
- oxygen sensor
- Kalman filter
- state machine

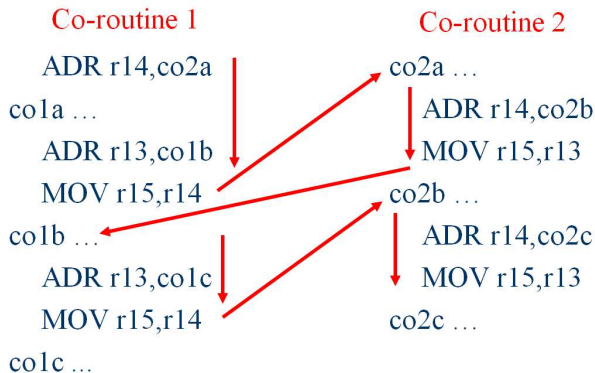
Typical Rates in Engine Control

Variable	Full range time (ms)	Update period (ms)
Engine spark timing	300	2
Throttle	40	2
Air flow	30	4
Battery voltage	80	4
Fuel flow	250	10
Recycled exhaust gas	500	25
Status switches	100	20
Air temperature	Seconds	400
Barometric pressure	Seconds	1000
Spark (dwell)	10	1
Fuel adjustment	80	8
Carburetor	500	25
Mode actuators	100	100

- Code turns into a mess
 - interruptions of one task for another
 - spaghetti code
- Although it can still be developed under a strong discipline (e. g. round robin architecture)



- Commonly used in the early days of embedded computing to handle multiple processes without processes
- Useful example of the complexities of not using processes for certain applications



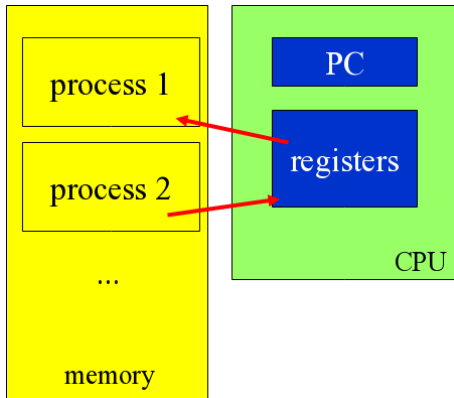
- Like subroutine, but caller determines the return address
- Co-routines voluntarily give up control to other co-routines
- Pattern of control transfers is embedded in the code

- Does not do nearly enough to help us construct complex programs with significant timing properties

- Can be ugly to work with
 - it's hard to trace through the flow of control

- Processes allow us to handle complex systems
 - fundamental abstraction for dealing with multiple simultaneous operations
- A process is a unique execution of a program
 - defined by its data and code
 - several copies of a program may run simultaneously or at different times
- A process has its own state
 - registers
 - memory
- The operating system manages processes

- **Activation record:** copy of process state
- Context switch
 - current CPU context goes out
 - new CPU context goes in



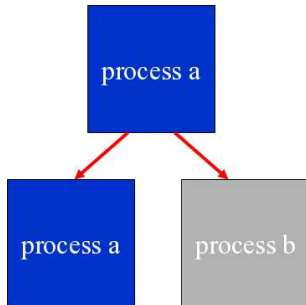
Thread (lightweight process)

- a process that shares memory space with other processes

Reentrancy

- ability of a program to be executed several times with the same results
- programs are easier to debug if they are reentrant

- Create a process with fork
 - parent process keeps executing old program
 - child process executes new program



- The fork process creates child:

```
childid = fork();
if (childid == 0) {
    /* child operations */
} else {
    /* parent operations */
}
```


- Overlays child code:

```
childid = fork();  
if (childid == 0) {  
    execv("mychild",childargs);  
    perror("execv");  
    exit(1);  
}
```

- It's the mechanism for moving the CPU from one executing process to another
- A process should not be able to tell that it was stopped and then restarted

Questions

- Who controls when the context is switched?
- How is the context switched?

- Improvement on co-routines
 - hides context switching mechanism
 - still relies on processes to give up CPU
- Each process allows a context switch call
- Separate scheduler chooses which process runs next
- Implementable with a function-queue architecture

Programming errors can keep other processes out

- Process never gives up CPU

- Process waits too long to switch, missing input

- Must copy all registers to activation record, keeping proper return value for PC

- Must copy new activation record into CPU state

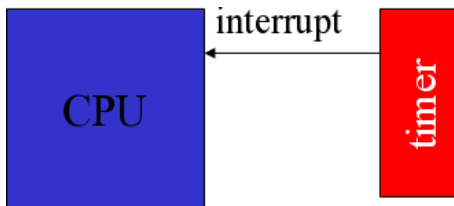
- Save old process:

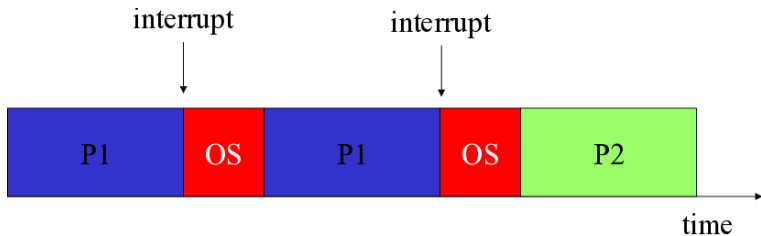
```
STMIA r13,{r0-r14}^ ; save registers
MRS r0,SPSR
STMDB r13,{r0,r15} ; save status register and PC
```

- Start new process:

```
ADR r0,NEXTPROC ; get address of next process
LDR r13,[r0]
LDMDB r13,{r0,r14}
MSR SPSR,r0 ; set status register
LDMIA r13,{r0-r14}^ ; get registers
MOVS pc,r14 ; restore PC
```

- Most powerful form of multitasking
- OS controls when contexts switches
- OS determines what process runs next
- Uses timer to call OS and to switch contexts





- Timer interrupt gives control to OS, which saves interrupted process's state in an activation record
- OS chooses next process to run
- OS installs desired activation record as current CPU state

How

- We could change the interrupt vector at every period
- When there was an interrupt the corresponding process would run

But

- We would need management code anyway
- We would have to know the next period's process at the start of the current process

- May want to test
 - context switch time assumptions
 - scheduling policy

- Can use OS simulator
 - to exercise process set
 - trace system behavior

- Processes can cause additional caching problems
 - even if individual processes are well-behaved, processes may interfere with each other

- Worst-case execution time with bad cache behavior is usually much worse than execution time with good cache behavior

- There are lots of operating systems targeted at embedded systems
- Used in a myriad of applications
- With more or less features
- Ported to great number of microprocessors

- Started as a collaboration between the University of California, Berkeley in co-operation with Intel Research and Crossbow Technology
- Targeted at Wireless Sensor Networks
- Programmed in nesC
- Interfaces and components for common abstractions
 - packet communication, routing, sensing, actuation and storage

- embedded real-time operating system
- 100% compatible with POSIX
 - capabilities for threads, communication, synchronization and I/O
- Programmed in C
- Targeted at high performance real time digital signal processing

- Real-time operating system
- Multitasking kernel with preemptive scheduling, fast interrupt response, memory management, interthread communication, mutual exclusion, event notification, and thread synchronization features
- Programmed in C
- Used in printer products of HP
 - targeted also at consumer electronics, medical devices, data networking applications, and SoC development

- Open-source Unix-like operating system
- Derived from the BSD
- Ported to a large number architectures
 - "Of course it runs NetBSD"
- Used in routers, switchers and "a toaster"



- Real-time operating system
- Designed to be small and simple
 - only three C source files
- Provides methods for multiple threads or tasks, mutexes, semaphores and software timers
- Good portability

- Pronounced "you-see-Linux"
- Fork of the Linux kernel for microcontrollers
 - without a memory management unit
- Used in network routers, security cameras, DVD/MP3 players, VoIP phone or Gateways, scanners and card readers.

- Provides multitasking and built-in TCP/IP stack
- Run on devices that are severely constrained in terms of memory, power, processing power, and communication bandwidth
 - focus on low-power wireless internet of things devices
- Used in street lighting systems, sound monitoring for smart cities, radiation monitoring systems, and alarm systems

- Smartphones
 - Android, iOS, ...
- Cisco IOS
- brickOS, leJOS
- many, many more

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- The process abstraction
- Context switching
- Multitasking

Computers as Components: Principles of Embedded Computing System Design , Marilyn Wolf. Morgan Kaufman. Ch. 6.1-6.4

- Scheduling policies